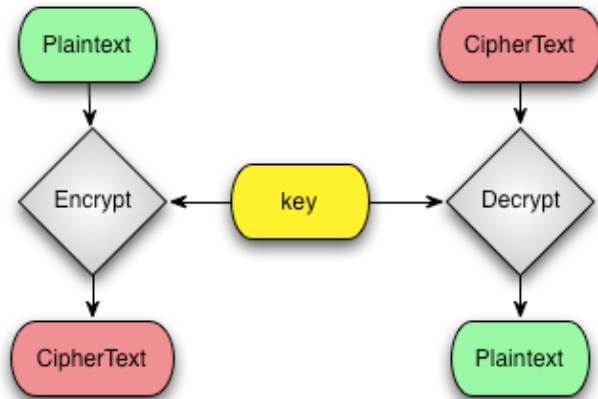


PKI, X.509 CERTIFICATES AND HTTPS WEBSERVERS

Public Key algorithms, digital certificates and PKI

PRELIMINARIES

Symmetric/Asymmetric cryptography



symmetric

The encryption and decryption keys are the same or can be directly derived from each other. Both keys are kept secret.

Examples: 3DES, AES, Blowfish, RC4

asymmetric

Encryption/decryption keys are different and it is **computationally unfeasible** to derive them from each other.

The encryption key be distributed, the other has to be kept secret.

For this reason it is also called Public Key cryptography.

Examples: RSA, Diffie-Hellman, ElGamal



Public Key cryptography: encryption/decryption

Alice



Alice wants to send a message
 M encrypted for Bob

Bob



Gets Bob's public key B_{pub}
(Somehow) verifies B_{pub} authenticity
Encrypts M with $B_{pub} \rightarrow C = F(B_{pub}, M)$

Alice sends C to Bob



Decrypts C with Bob's private key B_{priv}
 $M = F(B_{priv}, C)$

Note:

- 1) Only Bob can decrypt C
- 2) Nobody "can" derive B_{priv} from B_{pub}
- 3) This procedure can be inverted to implement a **digital signature**

Public Key cryptography: digital signature

Alice



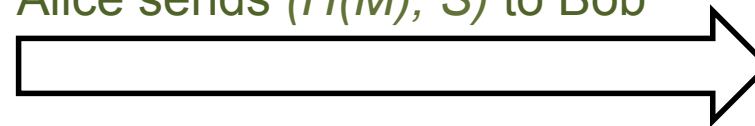
Alice wants to sign a message M so that Bob can verify its authenticity

Bob



Gets his own private key A_{priv}
Computes a hash of the message $H(M)$
Signs $H(M)$ with $A_{priv} \rightarrow S = F(A_{priv}, H(M))$

Alice sends $(H(M), S)$ to Bob



Computes a hash of the message $H(M)$
Verify the signature by verifying the following:
 $H(M) = F(A_{pub}, H(M)) ?$

Note:

- 1) Only Alice can sign M
- 2) Nobody can modify M and compute a valid signature S without knowing A_{priv}
- 3) Alice can include a nonce (given by Bob) in the signature to avoid a third entity to reuse the same signature for the same message M

RSA: key generation

1. Extract two “big” prime numbers p e q (**random, secret**)
2. Compute the RSA modulus: $N = p \times q$
3. Compute $\Phi(N) = (p - 1)(q - 1)$ (Eulero’s function)
4. Randomly generates the the number e : $1 < e < \Phi(N)$ relatively prime to $\Phi(N)$
5. Compute the number d : $e \times d = 1 \text{ mod } \Phi(N)$, or in other words e is the inverse of d in the group $\Phi(N)$

PUBLIC KEY: (N, e)

PRIVATE KEY: d

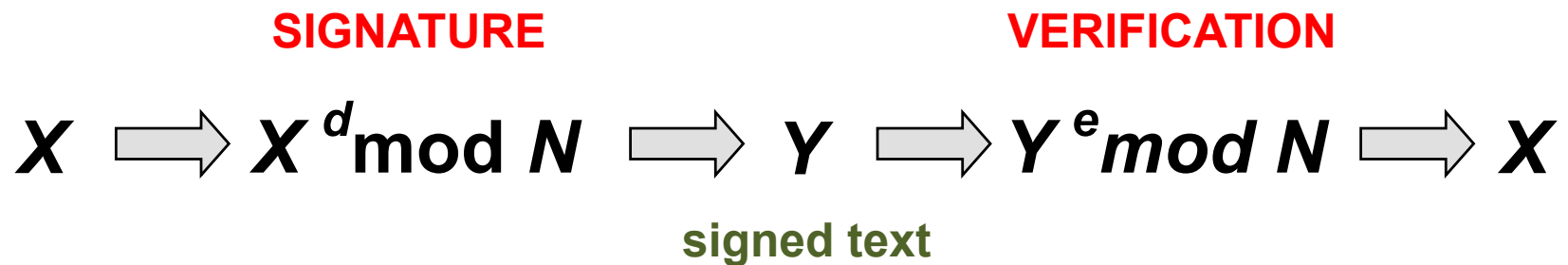
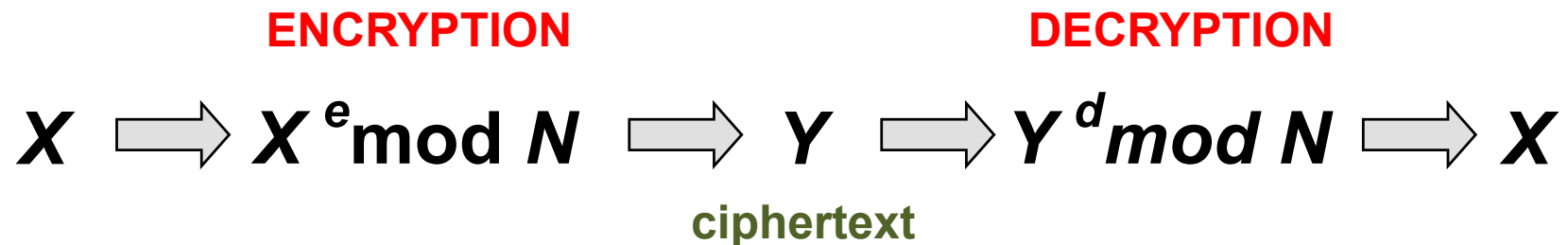
Must be kept secret: $p, q, \Phi(N), d$

Note:

- 1) to derive d from e an attacker should compute e^{-1} in $\Phi(N)$
- 2) $\Phi(N)$ is the number of integers less than or equal to n that are relatively prime to N
 - 2.1) to compute $\Phi(N)$ an attacker should know p and q (otherwise it’s unfeasible)
- 3) it is computationally unfeasible to factorize the product of two “big enough” prime numbers

RSA transformations

RSA transformation is simply a modular exponentiation with respectively the public private key



RSA with mathematica

```
Running...Untitled-1

In[1]:= p = RandomPrime[2^512]
Out[1]= 2 342 430 759 282 247 717 650 530 181 442 040 148 673 740 046 224 514 240 856 555 364 140 631 652 042 275 569 677 372 175 777 643 111 127 433 662 545 776 110 566 305 759 926 789 973 381 458 462 356 406 \
213

In[2]:= q = RandomPrime[2^512]
Out[2]= 1 139 470 330 073 101 760 388 131 330 594 903 773 840 813 370 432 508 808 717 960 489 897 996 311 831 461 021 155 156 382 389 908 783 099 506 649 014 761 011 252 607 146 489 100 588 660 943 874 346 780 \
073

In[3]:= modulus = p*q
Out[3]= 2 669 130 350 452 729 182 063 835 862 865 607 735 978 023 901 060 441 442 342 416 090 624 307 559 409 655 892 362 353 654 564 779 567 445 360 232 206 556 605 069 861 836 487 142 629 732 068 612 936 756 \
150 502 811 943 955 978 647 301 496 045 479 557 387 455 394 037 459 573 774 312 059 788 171 027 814 767 026 185 433 055 897 896 554 087 437 210 307 755 703 337 677 522 979 849 969 757 560 311 460 561 \
793 549

In[4]:= fi = (p - 1) * (q - 1)
Out[4]= 2 669 130 350 452 729 182 063 835 862 865 607 735 978 023 901 060 441 442 342 416 090 624 307 559 409 655 892 362 353 654 564 779 567 445 360 232 206 556 605 069 861 836 487 142 629 732 068 612 936 756 \
147 020 910 854 600 629 169 262 834 533 442 613 464 940 840 620 802 550 724 737 543 934 132 399 850 893 289 594 600 527 339 729 002 193 210 269 996 195 166 215 858 610 073 434 079 195 517 909 123 858 \
607 264

In[5]:= e = RandomPrime[fi]
Out[5]= 2 497 866 473 138 812 645 967 318 818 461 569 127 845 371 761 779 831 255 436 577 553 432 312 858 979 912 753 436 202 297 826 309 194 174 838 999 837 833 475 322 744 535 994 579 071 358 422 568 471 386 \
055 193 496 782 849 807 469 672 565 993 023 048 820 741 289 696 023 243 036 651 937 329 763 085 482 734 879 427 987 148 237 524 646 468 071 545 495 884 396 228 570 332 929 234 004 825 581 021 576 545 \
512 919

In[6]:= d = PowerMod[e, -1, fi] (*d=e^-(1)mod fi*)
Out[6]= 2 484 969 694 028 556 821 477 264 583 951 858 707 791 154 570 202 655 578 547 129 182 001 696 777 070 080 395 759 331 783 228 045 776 870 244 264 768 084 983 417 145 524 185 119 541 586 052 748 536 186 \
167 201 509 544 899 265 552 472 077 038 073 309 382 418 373 342 132 620 473 143 977 354 844 412 628 875 534 029 334 182 601 183 175 087 298 870 183 082 118 475 580 169 090 229 003 899 344 327 861 287 \
919 847

In[7]:= plain = 1667 850 607 (*ciao in ASCII*)
Out[7]= 1 667 850 607

In[8]:= ctext = PowerMod[plain, e, modulus] (*"ciao"^(e) mod N*)
Out[8]= 9 224 434 998 406 953 366 807 700 047 504 982 848 322 071 666 583 934 383 443 571 567 024 586 729 727 798 269 577 919 685 285 315 044 505 291 888 826 840 599 687 014 292 035 599 697 290 536 071 179 679 \
527 783 660 439 233 687 557 184 020 635 918 156 792 381 744 158 925 620 447 423 502 375 549 530 296 383 535 404 343 526 925 321 058 097 632 806 650 379 230 929 687 424 573 796 837 214 680 877 217 003 \
295 496

In[9]:= (*decifriamo elevando alla d*)
res = PowerMod[ctext, d, modulus]
Out[9]= 1 667 850 607

In[10]:= fi = EulerPhi[modulus] (*non ce la farà,dovrò abortire*)
Out[10]= $Aborted
```



PID	Process Name	User	% CPU	Threads	Real
61201	MathKernel	marlon	100.1	16	:
2994	VLC	marlon	4.8	14	1:
30629	Google Drive	marlon	1.2	15	:

...not exactly the real algorithm, but the concepts are the same!

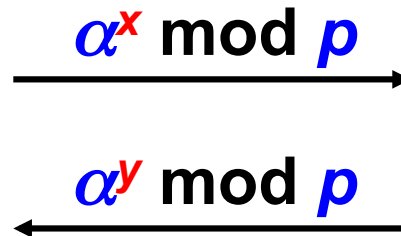
Diffie-Hellman Key exchange algorithm

Public: α, p

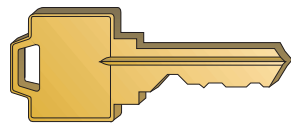
Secret: x, y

GOAL: exchange a common secret that only Alice and Bob can derive

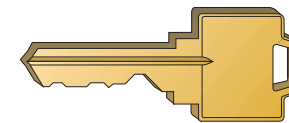
Random x



Random y



COMMON KEY



$$K = (\alpha^y)^x \bmod p$$

$$K = (\alpha^x)^y \bmod p$$

Note:

- 1) Common secret number exchanged with an asymmetric algorithm
- 2) to compute K from $(\alpha^x \bmod p)$ and $(\alpha^y \bmod p)$ an attacker should be able to compute the discrete logarithm $x = \log_{\alpha}(\alpha^x \bmod p)$ and $y = \log_{\alpha}(\alpha^y \bmod p)$...
- 3) ...which is computationally unfeasible for an attacker with "limited computational resources"

How does Alice obtain Bob's public key?

- Everything's perfect, you believe that nobody can break the public key algorithms if the numbers are "big enough"
- How are the public keys distributed?
 - In a network with n nodes, $n(n-1)/2$ keys have to be distributed!
 - What if my private key is lost or stolen? Should I need to notify all the remaining $(n-1)$ nodes to revoke my public key?
 - **Solution:** centralized or opportunistic distribution! (obvious, the public key don't have to be kept secret!)
- OK, the scalability issue is solved, but how can I be sure that a public key is authentic? How can Alice get the public key of Bob and be sure that it's really his?
- **SOLUTION:**
 - A **trusted** third party that issues some kind of proof that a public key is really related to a given identity

Public Key Certificate

- A public key certificate is a data structure that binds a public key (and therefore the related private key) to the identity of the legitimate owner $\rightarrow \text{CERT}_{ID}:\{\text{ID}, \text{Pub}_{ID}\}$
- The binding between $\{\text{ID}, \text{Pub}_{ID}\}$ is granted by a trusted certification authority that signs CERT_{ID}
- Provided that we have the CA's public key, we can verify the CA signature and therefore verify the public key authenticity

EXAMPLE:

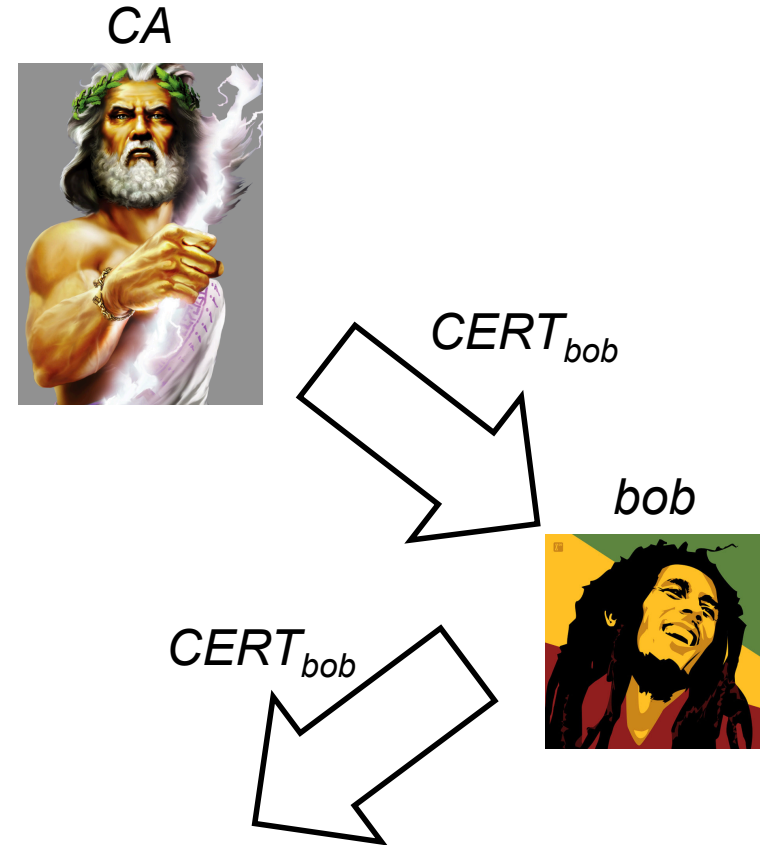
CA issues a public certificate for bob CERT_{bob}

CERT_{bob} contains:

- 1) Pub_{bob}
- 2) CA identity CA_{id}
- 3) CA signature of CERT_{bob}

Once I have the authentic Pub_{bob} , I just need to verify that the party I'm communicating with is actually Bob (i.e.: it has the private key)

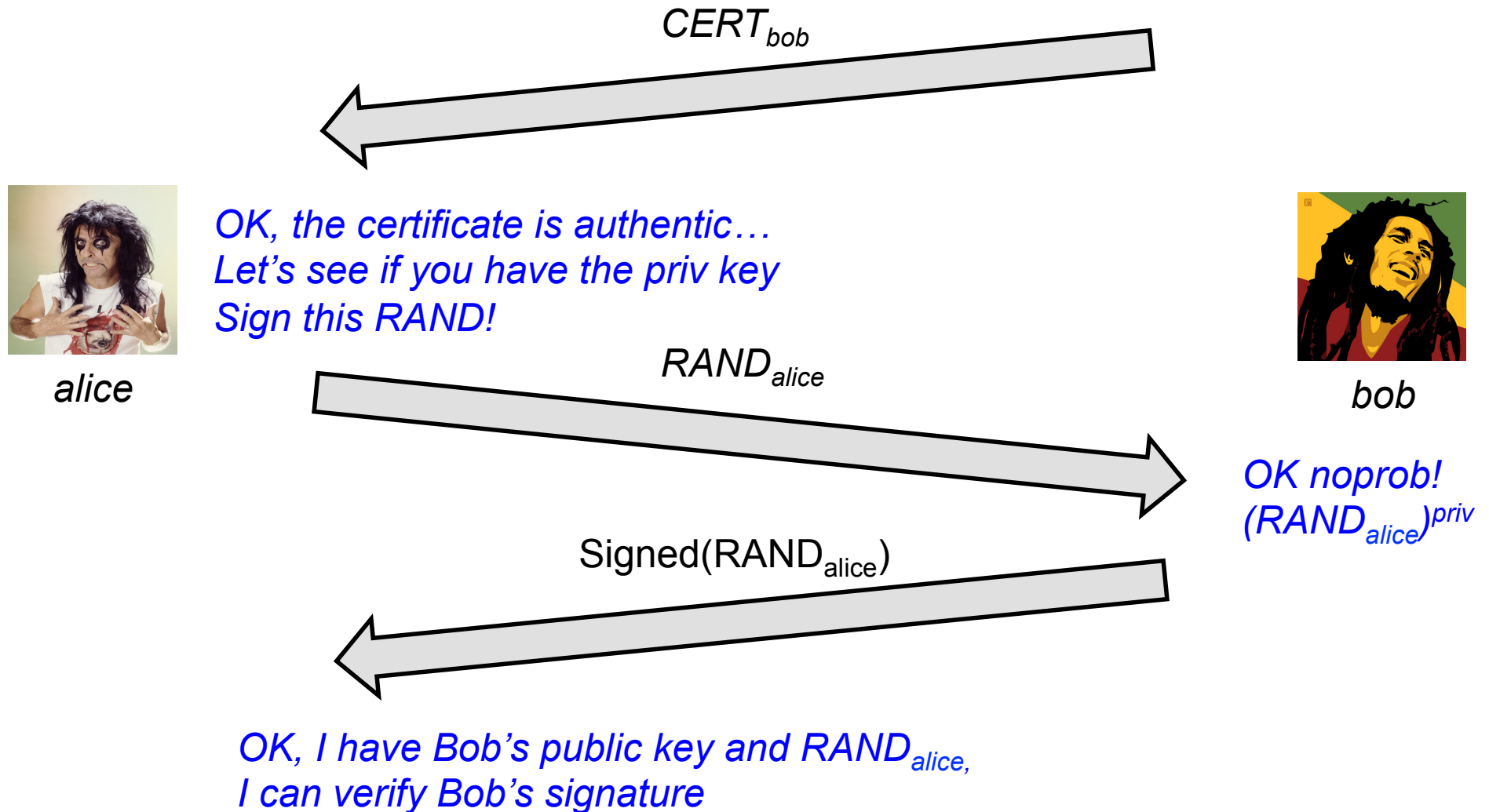
To do so, I perform a simple challenge/response mechanism. I extract a nonce and challenge Bob to sign this random number. Since the public key is authentic, and Bob couldn't know the random number, only the real Bob can sign the nonce correctly (and I can verify it)



alice

- I trust CA and I have CA's public key
- Verify CA signature $\text{CERT}_{bob} \rightarrow \text{OK!}$
- Pub_{bob} is authentic
- I can encrypt a message for Bob

Challenge/Response concept



Public Key Infrastructure

- A PKI consists of the protocols, the policies and the cryptographic mechanism used to manage the management of public key certificate
 - Creation, distribution, revocation, etc...
- A PKI requires the definition of:
 - Certificate format
 - Relationship among CAs
 - Mechanisms and policies for issuing and revoking certificate
 - Storage services
- Typical certificate format: X.509

X.509 format (high level)

Version, Validity, Serial Number, and others..
CA Identity
Subject Identity
Subject Public Key
CA Signature

X.509 certificate: real example



```
marlon@ubuntu:~/Desktop$ openssl x509 -in www.facebook.com
-----BEGIN CERTIFICATE-----
MIIGmjCCBRqgAwIBAgIQDG/IWVf6H1/JZyyf51zb5jANBgkqhkiG9w0BAQUFADBm
MQswCQYDVQQGEwJVUzEVMBMGA1UEChMMRG1naUNlcnQgSW5jMRkwFwYDVQQLExB3
d3cuZGlnaWNoLnQuY29tMSUwIiwvY29tMSUwIiwvY29tMSUwIiwvY29tMSUwIiwv
ZSBDOz0zMB4XDTEwMTEwMjYyMjYyMjYyMjYyMjYyMjYyMjYyMjYyMjYyMjYyMjYy
BhMCVVMxZzARBGNVBAgTCKNhbG1mb3JuaWEeEjAQBGNVBACTCVBhbG8gQWx0bzEX
MBUGA1UEChMORmFjZWJvb2s5IEluYy4xGTAXBGNVBAMTEHd3dy5mYWNlYm9vay5j
b20wgZ8wDQYJKoZIhvcNAQEBBQADgY0AMIGJAoGBAMHffwNBvcTk+mUzE3jVYjeW
p2HzsZa/I466h6ftB/neLeuox7ytd6ZeJQMDNuNN99Dxq2byty4zFr4mD11BFn//
twCe+g6ZFWxSGtcKxq375ACiP9sEplZppe3Wh7aIxYP16Maz/8AOH52jhXDtonYU
e3A+77BCCzjWggAj3WN1AgMBAAGjggNaMIIDVjAfbGNVHSMEGDAWgBRQ6n0J2yn7
EI+e5QEg1N55mUiD9zAdBgNVHQ4EFQQUldKM7bs1W6BE6Y2XvR7Q1jjzj0QwKQYD
VR0RBCIwIIQd3d3LmZHY2Vib29rLmNvbYIMZmFjZWJvb2s5Y29tMHsGCCsGAQUF
BwEBBg8wbTAKBggrBgEFBQcwAYYYaHR0cDovL29jc3AuZGlnaWNoLnQuY29tMEUG
CCsGAQUBFzAChjlodHRwOi8vY29tMSUwIiwvY29tMSUwIiwvY29tMSUwIiwvY29t
aWdoQXNzdXJhbWNoeE9tMy5jcnQwDgYDVR0PAQH/BAQDAgWgMAWA1UdEwEB/wQC
MAAwZQYDVROBF4wXDAsoCqgKIYmaHR0cDovL2NybDMuZGlnaWNoLnQuY29tL2Nh
My0yMDEwaS5jcmwWLAQoCiGJmh0dHA6Ly9jcmw0LmRpZ21jZXJ0LmNvbS5jYTMt
MjAxMGkuY3J5MIIBxgYDVR0gBIIBvTCCAbkwggG1BgtghkgBhv1sAQMAATCCAaQw
OgYIKwYBBQUHAQEwLmh0dHA6Ly93d3cuZGlnaWNoLnQuY29tL3NzbC1jcmHtcmVw
b3NpdG9yeS5odG0wggFkBggrBgEFBQcCAjCCAVYegFSAEEAbG5ACAAdQBzAGUA
IABvAGYAIAB0AGgAaQBzACAQwBIAHIAAdABpAGYAaQBjAGEAdABIACAAYwBvAG4A
cwb0AGkAdAB1AHQAZQBzACAAYQBJAGMAZQBwAHQAYQBuAGMAZQAgAG8AZgAgAHQA
aAB1ACAARABpAGcAaQBDAQUAcgB0ACAQwBQAC8AQwBQAFMAIABHAG4AZAAgAHQA
aAB1ACAAYwB1AGwAeQBpAG4AZwAgAFAAYQByAHQAeQAgaEEAZwByAGUAZQBtAGUA
bgB0ACAAdwBoAGkAYwBoACAAbABpAG0AaQB0ACAAbABpAGeAYgBpAGwAaQB0AHKA
IABHAG4AZAAgAGEAcgB1ACAaQBuAGMAbWByAHAAbWByAGEAdAB1AGQAIAB0AGUA
cgB1AGkAbgAgAGIAeQAeAHIAZQBmAGUAacgB1AG4AYwB1AC4wHQYDVRO1BBYwFAYI
KwYBBQUHAwEGCCsGAQUFBwMCMA0GCSqGSIb3DQEBBQUAA4IBAQA1M16QP60C/t6S
0p4S9+8Wao26ZqBarmZ2vEoSE+0S1vcPllwB1SDo8P2sZt4kGK/uo9fo+sectYg
GtLgJwcNev+lj30Hf06ZgQBqjYcNjcaAFsUd2AY39+wD6KK0QfYvdQwUAdF1p1aY
8DggH3cVeau14wQKd8nDtZlXdk80bncaYTDvmrpTUT9RPPXAtdMQgl+kmE0DDGeRB
2Sb30UvyoaTdtQXFvUJlhcsppgGHW14e6yCX+hXG70mZjUkkLHWqAYkM8J/w8Khwu
gqeCEJjrS1oyfLGPXDkAx9xtb3+v2DdAE0j8xCWg/hvleSrYh1SBXmU1zHyHHVE
yie0b6nD
-----END CERTIFICATE-----
```

Certificate Viewer: www.facebook.com

General Details

This certificate has been verified for the following usages:
SSL Server Certificate

Issued To

Common Name (CN)	www.facebook.com
Organization (O)	Facebook, Inc.
Organizational Unit (OU)	<Not Part Of Certificate>
Serial Number	0C:6F:C8:59:57:FA:1F:5F:C9:67:2C:9F:E6:5C:DB:E6

Issued By

Common Name (CN)	DigiCert High Assurance CA-3
Organization (O)	DigiCert Inc
Organizational Unit (OU)	www.digicert.com

Validity Period

Issued On	11/15/10
Expires On	12/3/13

Fingerprints

SHA-256 Fingerprint	BB A9 12 B4 FE 2F 26 88 7D 79 0B C4 2F 7A 98 7B C8 D8 1C 21 B1 90 C4 46 5B C3 1A 2C 5B 6F D2 31
SHA-1 Fingerprint	63 08 84 E2 79 CB 11 07 F1 FB 8A 6B 11 A6 4D 1B 14 76 3F 8E

Close

X.509 certificate: real example

Version: 3 (0x2)

Serial Number:

0c:6f:c8:59:57:fa:1f:5f:c9:67:2c:9f:e6:5c:db:e6

Signature Algorithm: sha1WithRSAEncryption

Issuer: C=US, O=DigiCert Inc, OU=www.digicert.com, CN=DigiCert High Assurance CA-3

Validity

Not Before: Nov 15 00:00:00 2010 GMT

Not After : Dec 2 23:59:59 2013 GMT

Subject: C=US, ST=California, L=Palo Alto, O=Facebook, Inc., CN=www.facebook.com

Subject Public Key Info:

Public Key Algorithm: rsaEncryption

RSA Public Key: (1024 bit)

Modulus (1024 bit):

00:c1:df:7d:63:41:bd:c4:e4:fa:65:33:13:78:d5: (... cut...) 0b:38:d6:82:00:23:dd:63:75

Exponent: 65537 (0x10001)

X509v3 extensions: (cut)

X509v3 Subject Key Identifier:

AA:57:4A:33:B6:EC:D5:6E:81:13:A6:36:5E:F4:7B:43:58:F3:8F:44

X509v3 Subject Alternative Name:

DNS:www.facebook.com, DNS:facebook.com

X509v3 Key Usage: critical

Digital Signature, Key Encipherment

X509v3 Basic Constraints: critical

CA:FALSE

X509v3 Extended Key Usage:

TLS Web Server Authentication, TLS Web Client Authentication

Signature Algorithm: sha1WithRSAEncryption

25:33:5e:90:3f:ad:02:fe:de:92:d2:9e:12:f7:ef:16:6a:8d: (... cut...) 8e:6f:a9:c3

Certificate Signing Request

- A certificate signing request (also CSR or certification request) is a message sent from an applicant to a certificate authority in order to apply for a digital identity certificate
- The most common format for CSRs is the PKCS#10 specification
- Operations:
 - the applicant first generates a key pair, keeping the private key secret
 - the applicant generates a CSR contains information identifying herself (X.509 subject field), optional X.509 extensions (e.g. key usage: RSA authentication for web servers) and the public key chosen by the applicant
 - The CSR may be accompanied by other credentials or proofs of identity required by the certificate authority, and the certificate authority may contact the applicant for further information

Let's build our own certification authority

OPENSSL X509 TUTORIAL

OpenSSL

- OpenSSL is a cryptography toolkit implementing the Secure Sockets Layer (SSL v2/v3) and Transport Layer Security (TLS v1) network protocols and related cryptography standards required by them
 - www.openssl.org
- Main component
 - Cryptography library: `libcrypto`
 - SSL/TLS protocol library: `libssl`
 - `openssl` program
- The `openssl` program is a command line tool for using the various cryptography functions of OpenSSL's crypto library from the shell. It can be used for
 - Creation and management of private keys, public keys and parameters
 - Public key cryptographic operations
 - Creation of X.509 certificates, CSRs and CRLs
 - Calculation of Message Digests
 - Encryption and Decryption with Ciphers
 - SSL/TLS Client and Server Tests
 - Handling of S/MIME signed or encrypted mail
 - Time Stamp requests, generation and verification

Create a CA and sign certificate request with openssl

- Typical workflow
 1. Generate the RSA key pair for our CA
 2. Create a self-signed certificate for our CA
 3. Generate the RSA key pair for the web server
 4. Generate a CSR for the web server
 5. Sign the CSR with the CA private key

- Very simple Lab-pki
 - Create the CA and issue the certificates (single level certification ROOT_CA→certificate) with openssl from the host machine
 - Create a netikit lab (Lab9-pki) with just one VM (with a TAP 10.0.0.1,10.0.0.2) that will be our test web server
 - Setup Apache2 for a HTTPS website

Create the CA keys

Prepare our CA folder and the serial number file

```
marlon@marlon-vmxnb:~/Labs$ mkdir cgr1CA
marlon@marlon-vmxnb:~/Labs$ cd cgr1CA/
marlon@marlon-vmxnb:~/Labs/cgr1CA$ echo -e "01\n" > serial
```

Create the CA key pair

```
marlon@marlon-vmxnb:~/Labs/cgr1CA$ openssl genrsa -out ca.key 2048
Generating RSA private key, 2048 bit long modulus
.....+++
.....+++
e is 65537 (0x10001)
```

Generate the CA self signed certificate

This command will create a self signed certificate, i.e. a certificate where the issuer and the subject are the same entities

```
marlon@marlon-vmxbn:~/Labs/cgriCA$ openssl req -new -x509 -days
3650 -key ca.key -out ca.crt
You are about to be asked to enter information that will be
incorporated
into your certificate request.
What you are about to enter is what is called a Distinguished Name
or a DN.
There are quite a few fields but you can leave some blank
For some fields there will be a default value,
If you enter '.', the field will be left blank.
-----
Country Name (2 letter code) [AU]:IT
State or Province Name (full name) [Some-State]:
Locality Name (eg, city) []:Rome
Organization Name (eg, company) [Internet Widgits Pty Ltd]:cgriCA
Organizational Unit Name (eg, section) []:
Common Name (eg, YOUR name) []:cgri-cert-authority
Email Address []:ca@cgri.edu
```

Let's take a look at our first certificate

```
marlon@marlon-vmxbn:~/Labs/cgriCA$ openssl x509 -in ca.crt -text -noout
Certificate:
  Data:
    Version: 3 (0x2)
    Serial Number:
      b6:ef:85:6f:71:e5:68:bb
    Signature Algorithm: sha1WithRSAEncryption
    Issuer: C=IT, ST=Some-State, L=Rome, O=cgriCA, CN=cgri-cert-authority
                                                    emailAddress=ca@cgri.edu
  Validity
    Not Before: May 24 10:44:00 2012 GMT
    Not After : May 22 10:44:00 2022 GMT
  Subject: C=IT, ST=Some-State, L=Rome, O=cgriCA, CN=cgri-cert-authority/
                                                    emailAddress=ca@cgri.edu
  Subject Public Key Info:
    Public Key Algorithm: rsaEncryption
      Public-Key: (2048 bit)
      Modulus:
        00:a1:2c:f1:bf:a2:af:4a:3a:6e:f7:e7:13:b5:42:
        32:4c:2c:d2:3b:0f:09:68:d6:67:6e:af:05:23:a8:
        59:eb:ef:85:19:7c:75:18: Cut!
```

Let's make the web server keys and CSR

Create the subject's (i.e. our web server) key pair

```
marlon@marlon-vmxbn:~/Labs/cgriCA$ openssl genrsa -out server.key 1024
Generating RSA private key, 1024 bit long modulus
.++++++
.....++++++
e is 65537 (0x10001)
```

Create the subject's CSR. This certificate will be signed with the CA's private key

```
marlon@marlon-vmxbn:~/Labs/cgriCA$ openssl req -new -key server.key -out
server.csr

Country Name (2 letter code) [AU]:IT
State or Province Name (full name) [Some-State]:
Locality Name (eg, city) []:Rome
Organization Name (eg, company) [Internet Widgits Pty Ltd]:
Organizational Unit Name (eg, section) []:
Common Name (eg, YOUR name) []:testssl.cgri.edu ←
Email Address []:testssl@cgri.edu
```

This has to be
The web site FQDN

CSR signing

This command will sign the CSR with the CA's private key

```
marlon@marlon-vmxnb:~/Labs/cgr1CA$ openssl x509 -req -in server.csr -out
server.crt -sha1 -CA ca.crt -CAkey ca.key -CAserial serial -days 3650
Signature ok
subject=/C=IT/ST=Some-State/L=Rome/O=Internet Widgits Pty Ltd/
CN=testssl.cgr1.edu/emailAddress=testssl@cgr1.edu
Getting CA Private Key
```

Dump the signed certificate

```
marlon@marlon-vmxnb:~/Labs/cgr1CA$ openssl x509 -in server.crt -text -noout
Certificate:
  Data:
    Version: 1 (0x0)
    Serial Number: 3 (0x3)
    Signature Algorithm: sha1WithRSAEncryption
    Issuer: C=IT, ST=Some-State, L=Rome, O=cgr1CA, CN=cgr1-cert-authority/
    emailAddress=ca@cgr1.edu
  Validity
    Not Before: May 24 10:50:25 2012 GMT
    Not After : May 22 10:50:25 2022 GMT
  Subject: C=IT, ST=Some-State, L=Rome, O=Internet Widgits Pty Ltd,
    CN=testssl.cgr1.edu/emailAddress=testssl@cgr1.edu
  Subject Public Key Info:
    Public Key Algorithm: rsaEncryption
```

How to protect our web server

HTTPS SERVER WITH APACHE2

Let's configure Apache2

We are going to create a virtual host for the website “testssl.cgrl.edu” in the netkit lab “Lab9-pki”

Configuration file, keys and certificate already in server:root/
Webserver media file and index.html in server:/var/www/testssl

Set-up everything properly before enabling the new site

- Configuration file testssl.cgrl.edu goes into /etc/apache2/site-available
- Keys and Certificate in the proper directory (see the conf file)

Run the following commands:

```
server# a2ensite testssl.cgrl.edu
```

Enable our HTTPS web site

```
server# a2enmod ssl
```

Enable Apache2 SSL module

```
server# /etc/init.d/apache2 start
```

Start Apache2
(or “restart” if already up)

testssl.cgri.edu config file

```
IfModule mod_ssl.c>
<VirtualHost _default_:443>
DocumentRoot "/var/www/testssl"

ServerName testssl.cgri.edu:443
ServerAdmin testssl@cgri.edu

SSLEngine On
SSLCipherSuite HIGH:MEDIUM
SSLProtocol all -SSLv2
SSLCertificateFile /etc/apache2/ssl/server.crt
SSLCertificateKeyFile /etc/apache2/ssl/server.key
SSLCertificateChainFile /etc/apache2/ssl/ca.crt
SSLCACertificateFile /etc/apache2/ssl/ca.crt

<Directory "/var/www/testssl">
    Options Indexes
    AllowOverride None
    Allow from from all
    Order allow,den
</Directory>
</VirtualHost>
</IfModule>
```

Connect to the server

Could not verify this certificate for unknown reasons.

Issued To	
Common Name (CN)	testssl.cgri.edu
Organization (O)	Internet Widgits Pty Ltd
Organizational Unit (OU)	<Not Part Of Certificate>
Serial Number	03

Issued By	
Common Name (CN)	cgri-cert-authority
Organization (O)	cgriCA
Organizational Unit (OU)	<Not Part Of Certificate>

Validity	
Issued On	05/24/2012
Expires On	05/22/2022

Fingerprints	
SHA1 Fingerprint	D2:FE:69:85:33:94:D8:56:DA:64:8B:DA:31:F8:63:E3:F0:01:26:6F:66:C1:94:80:2F:EE:62:F2:1
MD5 Fingerprint	

Technical Details

testssl.cgri.edu uses an invalid security certificate.

The certificate is not trusted because the issuer certificate is not trusted.

(Error code: sec_error_untrusted_issuer)

I Understand the Risks

If you understand what's going on, you can tell Firefox to start trusting this site's identification. **Even if you trust the site, this error could mean that someone is tampering with your connection.**

Don't add an exception unless you know there's a good reason why this site doesn't use trusted identification.

Unknown CA (of course...)

You can also manually and permanently add the certificate before trying to connect

Note: append the following line to the file `/etc/hosts` on the host machine
`testssl.cgri.edu 10.0.0.2`

TLsv1 trace with our certificate

The image shows a Wireshark 1.6.2 interface with a filter set to 'ssl'. The packet list table is as follows:

Time	Source	Destination	Protocol	Length	Info
4:0.000473	10.0.0.1	10.0.0.2	TLsv1	235	Client Hello
6:0.011022	10.0.0.2	10.0.0.1	TLsv1	1514	Server Hello
8:0.011170	10.0.0.2	10.0.0.1	TLsv1	844	Certificate, Server Key Exchange, Server Hello Done
10:0.013859	10.0.0.1	10.0.0.2	TLsv1	264	Client Key Exchange, Change Cipher Spec, Encrypted Handshake
11:0.019209	10.0.0.2	10.0.0.1	TLsv1	348	Encrypted Handshake Message, Change Cipher Spec, Encrypted
12:0.019530	10.0.0.1	10.0.0.2	TLsv1	439	Application Data
16:0.070438	10.0.0.2	10.0.0.1	TLsv1	678	Application Data, Application Data, Application Data, Appl
17:0.080485	10.0.0.1	10.0.0.2	TLsv1	455	Application Data

The details pane for the selected packet (8) shows the following structure:

- Certificates Length: 1750
- ▼ Certificates (1750 bytes)
 - Certificate Length: 769
 - ▶ Certificate (pkcs-9-at-emailAddress=testssl@cgrl.edu,id-at-commonName=testssl.cgrl.edu,id-at-organizationName=testssl.cgrl.edu) Certificate Length: 975
 - ▶ Certificate (pkcs-9-at-emailAddress=ca@cgrl.edu,id-at-commonName=cgrl-cert-authority,id-at-organizationName=cgrl-cert-authority) Certificate Length: 784
 - ▶ TLsv1 Record Layer: Handshake Protocol: Server Key Exchange
 - ▶ TLsv1 Record Layer: Handshake Protocol: Server Hello Done

Red arrows point to the 'issuer' field (top arrow) and the 'subject' field (bottom arrow) in the certificate details pane.

The hex dump at the bottom shows the following data:

```
0000 16 03 01 06 dd 0b 00 06 d9 00 06 d6 00 03 01 30 .....0
0010 82 02 fd 30 82 01 e5 02 01 03 30 0d 06 09 2a 86 ...0...0...*
0020 48 86 f7 0d 01 01 05 05 00 30 7c 31 0b 30 09 06 H.....011.0..
```

HTTP plaintext auth over SSL

- Safest way to authenticate via HTTP, better than digest auth
- You first create a secure channel with the authenticated web server
- You send authentication credentials in clear (from the HTTP point of view) but inside the secure (encrypted/authenticated) channel
- The test website already has the following password-protected directory

```
<Directory "/var/www/testssl/secret">
  AuthType Basic
  AuthName "Username and Password Required"
  AuthUserFile /etc/apache2/.htpasswd
  Require valid-user
</Directory>
```

To try it you need to grant access to a new user, for example: uid "007" password "jamesbond"

```
server# htpasswd -c -m /etc/httpd/.htpasswd 007
New password:
```

Client authentication via X509 certificate

- The client may authenticate itself with a X509 certificate
- To do so we need to
 1. Configure the web server to force SSL client authentication

```
<Directory "/var/www/testssl/cert-required">  
    SSLVerifyClient require  
    SSLVerifyDepth 1  
</Directory>
```

2. Create a client certificate and configure the web browser to use it (exported it in PKCS 12 format. **NOTE:** to use it with firefox you need to enable SSL renegotiation. With (my) chrome (v. 15.0.874.106 (Developer Build 107270 Linux) Ubuntu 11.10) it's already OK)

```
server# openssl genrsa -out client.key 1024  
server# openssl req -new -key client.key -out client.csr  
server# openssl x509 -req -in client.csr -out client.crt -sha1 -CA  
ca.crt -CAkey ca.key -CAserial serial -days 3650  
server# openssl pkcs12 -export -in client.crt -inkey client.key -  
out client.p12
```